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PROGRAM DESCRIPTION AND USER'S GUIDE FOR ZIPCAP -- A ZERO-ONE I--ETC(U)

JUL 80 K L CHABRA, R M SOLAND

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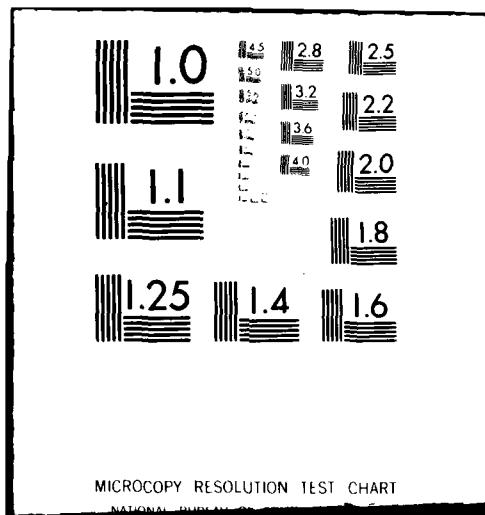


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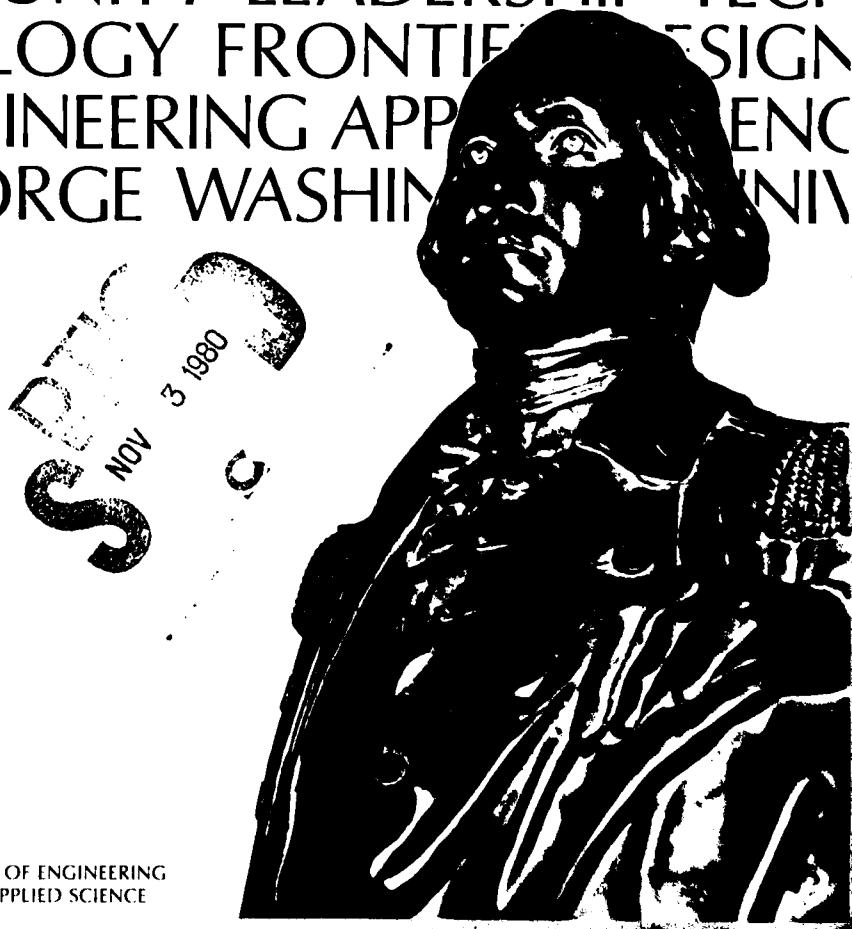
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PROGRAM DESCRIPTION AND USER'S GUIDE FOR ZIPCAP --  
A ZERO-ONE INTEGER PROGRAM TO SOLVE MULTIACTIVITY  
MULTIFACILITY CAPACITY-CONSTRAINED ASSIGNMENT  
PROBLEMS

by

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Richard M. Soland

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THE GEORGE WASHINGTON UNIVERSITY  
School of Engineering and Applied Science  
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Program in Logistics

Abstract  
of  
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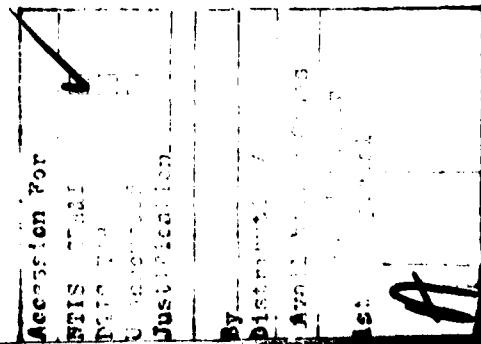
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1. Introduction

Multiactivity multifacility assignment problems are important because of their practical applications. The objective of such a problem is to minimize the total cost of the system, which includes both variable costs and fixed costs. A problem which has been of recent interest takes into account capacity constraints as well [Gross, Pinkus, and Soland (1979)]. A solution algorithm and a computer program implementing it have been developed for this kind of problem, i.e., for multiactivity multifacility 0-1 assignment problems with capacity constraints.

The development of the solution algorithm and computational test results using the computer program are described in a separate report titled "Solving a Multiactivity Multifacility Capacity-constrained 0-1 Assignment Problem."

This document, on the other hand, provides a description of the

computer program and instructions for its use. It is a FORTRAN program and is named ZIPCAP, an acronym for Zero-one Integer Program to solve multi-activity multifacility Capacity-constrained Assignment Problems.

Section 2 of this document provides a mathematical formulation of the basic problem and examples of application areas. Section 3 includes the computational steps and a program description, whereas Section 4 provides user information including program options and input data specifications. Test problems and output illustrations are covered in Section 5.

## 2. Problem Formulation

### 2.1 Basic Problem

The basic problem, called problem (P), is to find  $x_{ij}$  and  $y_k$  values that:

$$(P) \quad \left. \begin{array}{l} \text{Minimize} \quad \sum_{i=1}^m \sum_{j=1}^n a_{ij} x_{ij} + \sum_{k=1}^p b_k y_k \\ \text{subject to} \quad \sum_{i=1}^m x_{ij} = 1 \quad j=1, \dots, n \end{array} \right\} \quad (1)$$

$$\left. \begin{array}{l} \\ \sum_{i=1}^m \sum_{j=1}^n d_{ijk} x_{ij} \leq s_k y_k \quad k=1, \dots, p \end{array} \right\} \quad (2)$$

$$x_{ij}, y_k = 0 \text{ or } 1 \text{ for all } i, j, k \quad (3)$$

$$x_{ij}, y_k = 0 \text{ or } 1 \text{ for all } i, j, k \quad (4)$$

where  $m$ ,  $n$ , and  $p$  represent the number of designs, activities, and facilities respectively; and  $i$ ,  $j$ ,  $k$  are the corresponding indices.

The parameters are interpreted as follows:

$a_{ij}$  = variable cost of activity  $j$  using design  $i$ ,

$b_k$  = fixed cost of facility  $k$ ,

$s_k$  = capacity available at facility  $k$ , and

$d_{ijk}$  = capacity required at facility  $k$  for activity  $j$   
when activity  $j$  uses design  $i$ .

The decision variables are:

$x_{ij}$  = 1 if activity  $j$  uses design  $i$ ,  
= 0 otherwise.

$y_k$  = 1 if facility  $k$  is used,  
= 0 otherwise.

Simply stated, in problem (P) one seeks to minimize the sum of variable and fixed costs subject to constraint set (2) that each activity be assigned a single design and subject to constraint set (3) that the capacity constraints imposed by the facility capacities are satisfied.

Problem (P) has  $mn+p$  0-1 variables and  $n+p$  constraints.

## 2.2 Key Elements

The key elements in problem (P) are activities, facilities, designs, variable costs, and fixed costs. In general, these can be defined as follows.

- Activity -- a type of component, product or item under consideration.

- Facility -- an installation or location where activities can be served or serviced.
- Design -- a meaningful configuration of facilities along with a meaningful strategy for using them. For example, if there are four facilities 1, 2, 3 and 4, a design may include facilities 1, 2 and 4, whereas another design may include facility 2 only.
- Variable Cost -- the cost associated with a particular activity being assigned to a particular design.
- Fixed Cost -- the cost associated with a facility (such as part of its purchase, operation and maintenance cost) that is independent of the activities served at that facility.

### 2.3 Areas of Application

The computer program ZIPCAP is designed to solve a problem that can be formulated as problem (P). The formulation applies to existing as well as to proposed facilities. In other words, it is useful for a situation where the decision may be to delete some of the existing facilities, as well as for a situation involving a choice among proposed facilities (or locations).

Table 1 includes examples of areas where this formulation can be applied. Within each application area, activities and facilities are identified. The implications of designs, variable costs and fixed costs are apparent.

### 3. Computational Procedure and Program Description

The procedure to solve problem (P) basically involves branch and bound [Geoffrion and Marsten (1972)], and Lagrangean relaxation

Table 1EXAMPLES OF APPLICATION AREAS

<u>Area of Application</u>	<u>Activities</u>	<u>Facilities</u>
1. Design of multi-echelon inventory systems*	Items to be stocked	Warehouses (Comprising different levels or echelons, e.g., central, regional, and local warehouses)
2. Assignment of repairable components**	Major components of a unit, e.g., components of an aircraft, a ship, a piece of machinery	Repair depots
3. Design of training programs	Training program categories or occupational classifications	Training schools
4. Location of facilities***	Types of services, e.g., health-care services	Buildings or installations, e.g., health-care centers

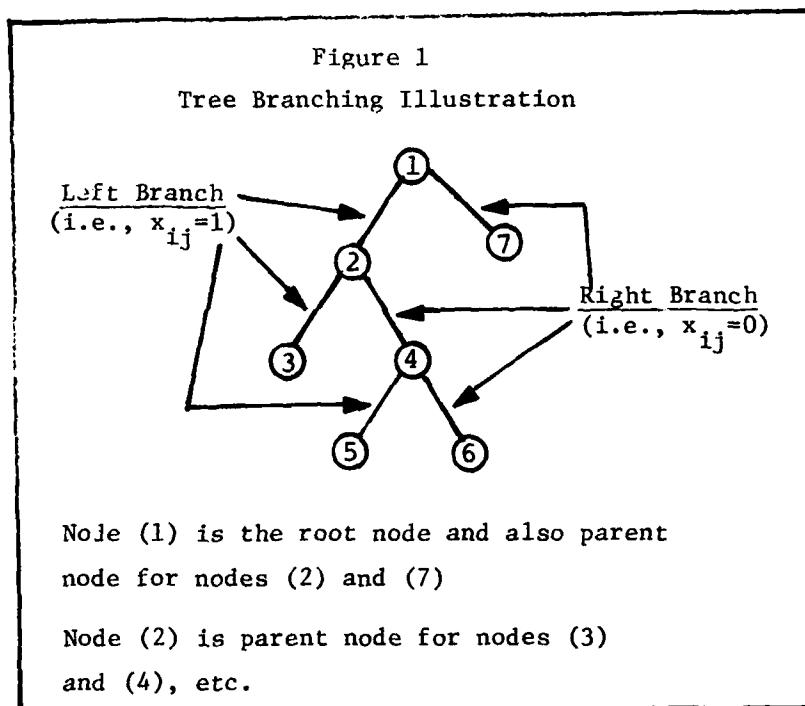
\* Gross, Pinkus, and Soland (1979)

\*\*Gross and Pinkus (1979)

\*\*\*Pinkus, Gross, and Soland (1973)

[Geoffrion (1974)]. In addition certain heuristics are used for branching, and to exclude infeasible assignments.

Branching is done on the  $x_{ij}$  variables by setting them equal to 1 or 0. The branching commences by setting a variable equal to 1 and moving to the left-branch node. When backtracking, the corresponding variable is set equal to 0 and we move to the right-branch node (if the right-branch has not yet been explored). Figure 1 illustrates a branching tree.



The  $x_{ij}$  variables which have been set equal to 1 or 0 at any given node are termed fixed variables, and the remaining ones are termed free variables.

### 3.1 Computational Steps

Figure 2 presents a simplified flow diagram, showing the

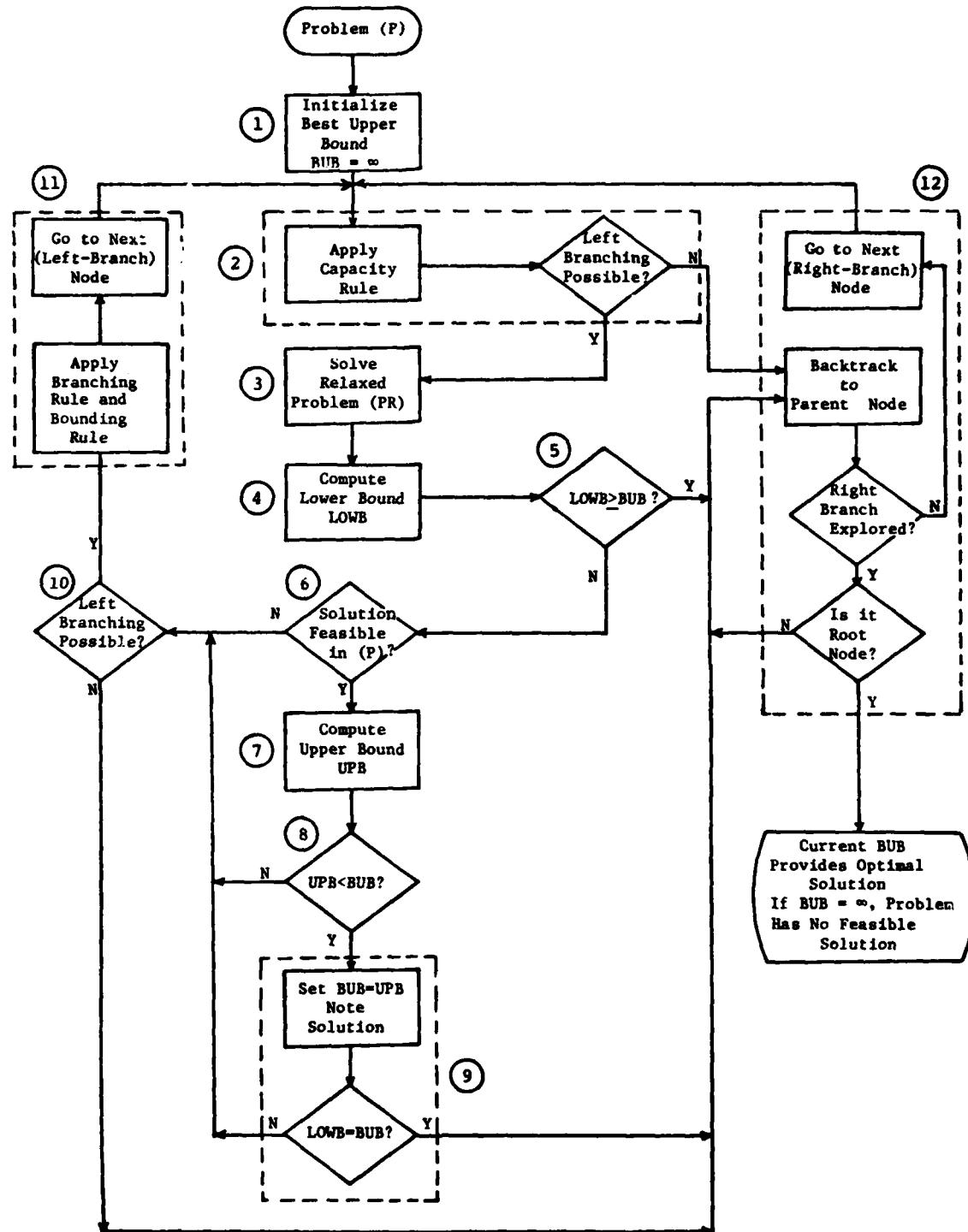


Figure 2  
Simplified Flow Diagram Showing Computational Steps

computational steps of the computer program. Following is a description of these steps.

Step 1. Initialize by setting best upper bound (BUB) equal to a large value and the node number to 1.

Step 2. Apply the capacity rule to exclude infeasible assignments prior to solving the relaxed problem.

Exclude a free  $x_{ij}$  if, for any facility  $k$  and activity  $j$ ,

$$\left( d_{ijk} - \min_j d_{ijk} \right) > \left( s_k - \sum_j \min_j d_{ijk} \right).$$

If an  $x_{ij}$  is already fixed as 1, the corresponding  $d_{ijk}$  replaces  $\min_j d_{ijk}$ .

If  $\sum_j \min_j d_{ijk} > s_k$ , then backtrack, i.e., go to step 12.

Step 3. Solve the relaxed problem (PR), i.e.,

$$(PR) \left\{ \begin{array}{l} \text{Minimize} \quad \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij} \\ \text{subject to} \quad \sum_{i=1}^m x_{ij} = 1 \quad j=1, \dots, n, \\ \qquad \qquad \qquad x_{ij} = 0 \text{ or } 1 \text{ for all } i, j, \end{array} \right.$$

$$\text{where } c_{ij} = a_{ij} + \sum_{k=1}^p b_k \left( 1 - \delta_k \right) \left( d_{ijk}/s_k \right),$$

$$\delta_k = 1 \text{ if } \sum_i \sum_j x_{ij} e_{ik} > 0,$$

such that  
 $x_{ij}$  is fixed

= 0 otherwise, and

$e_{ik} = 1$  if design  $i$  uses facility  $k$ ,

= 0 otherwise.

Step 4. Lower bound (LOWB) is the sum of the optimal solution value of problem (PR) and the fixed cost FC. FC is the cost of the facilities forced into the solution because of the  $x_{ij}$  variables fixed as 1:

$$FC = \sum_{k=1}^p \delta_k b_k$$

Step 5. Compare lower bound with best upper bound. If  $LOWB \geq BUB$ , go to step 12.

Step 6. Check if the solution of problem (PR) obtained in step 3 satisfies the capacity constraints of problem (P), i.e.,

$$\sum_{i=1}^m \sum_{j=1}^n d_{ijk} x_{ij} \leq s_k y_k \quad k=1, \dots, p,$$

where  $y_k = 1$  if  $\sum_i \sum_j d_{ijk} x_{ij} > 0$ ,

= 0 otherwise

If the capacity constraints are not satisfied, go to step 10.

Step 7. Compute upper bound (UPB), i.e.,

$$UPB = \sum_{i=1}^m \sum_{j=1}^n a_{ij} x_{ij} + \sum_{k=1}^p b_k y_k$$

Step 8. Compare upper bound with best upper bound. If UPB  $\geq$  BUB, go to step 10.

Step 9. Retain UPB as BUB, i.e., set BUB = UPB and note the corresponding solution comprising  $x_{ij}$  and  $y_k$  variables.

If LOWB equals BUB (same as UPB in this case), go to step 12.

Step 10 If an  $x_{ij}$  variable has been fixed as 1 for each of the n activities, then go to step 12.

Step 11. Select an  $x_{ij}$  variable for branching to the left node. The choice of the branching variable depends on the  $c_{ij}$  values and is such that the corresponding  $x_{ij}$ , if perturbed, has the maximum impact on the optimal value of problem (PR). That is, for each  $j$ , obtain  $D_j = c_{2j} - c_{1j}$  where  $c_{2j}$  is the second smallest permissible and  $c_{1j}$  is the smallest permissible  $c_{ij}$  value in column  $j$ .

A  $c_{ij}$  value is considered permissible if it does not correspond to an activity  $j$  having  $x_{ij}$  variable fixed as 1, or to an  $x_{ij}$  variable fixed as 0. The  $x_{ij}$  variables can be fixed as 1 or 0 because of the branching rule, backtracking, and/or use of the capacity rule. Further,  $x_{ij}$  may be fixed as 0 because of the bounding rule, i.e., if  $(c_{ij} - c_{1j}) > (BUB - LOWB)$ . The  $x_{ij}$  variable selected has the maximum value of  $D_j$  and has  $c_{ij} = c_{1j}$ .

Set the selected  $x_{ij}$  variable to 1. Increase the node number by one, and go to step 2.

Step 12. Backtracking continues until a node is reached for which

the right branch has not been explored. Set the corresponding  $x_{ij}$  variable to 0. Increase the node number by one, and go to step 2.

If no such node is found, it implies that the root node has been reached, and the procedure terminates.

The current BUB and the solution obtained in step 9 provide the optimal solution and its value. If BUB happens to be the initial value, the problem does not have a feasible solution, i.e., the capacity constraints cannot be satisfied.

### 3.2 Computer Program

The computer program ZIPCAP is written in FORTRAN IV. It is based on the flow diagram and the computational steps described in the preceding paragraphs.

Appendix A provides a listing of the program. Extensive use of comment cards provides details within each computational step.

A dictionary of the symbolic names used in the program listing is provided in Appendix B.

The program has been developed and tested on an IBM 3031 at The George Washington University.\* The only internally supplied subroutine used by the program is TIMET from the PIJ Library. This subroutine provides the elapsed time for program execution and excludes the time to read and write the input and output.

---

\* IBM 3031 CPU with 3 million Bytes Real Memory. Operating System OS/VSI Release 6.7. Various compilers including FORTRAN levels G 1 to 6 and H 1 to 6.

The program comprising about 480 lines is currently dimensioned for a problem size of 35 designs (m), 35 activities (n) and 30 facilities (p). The user capacity required to execute this program is 346 K bytes. The functional relationship between the dimensions of the arrays and the storage requirement is given by

$$f(m, n, p) = 4 \left[ (p+4)mn + (m+5)p + 9n \right] \text{ bytes.}$$

Thus the program size to execute a problem has two components: one, due to the program itself, comprises 173 K bytes, and the other, dependent on the dimensions of the arrays, is given by the above functional relationship.

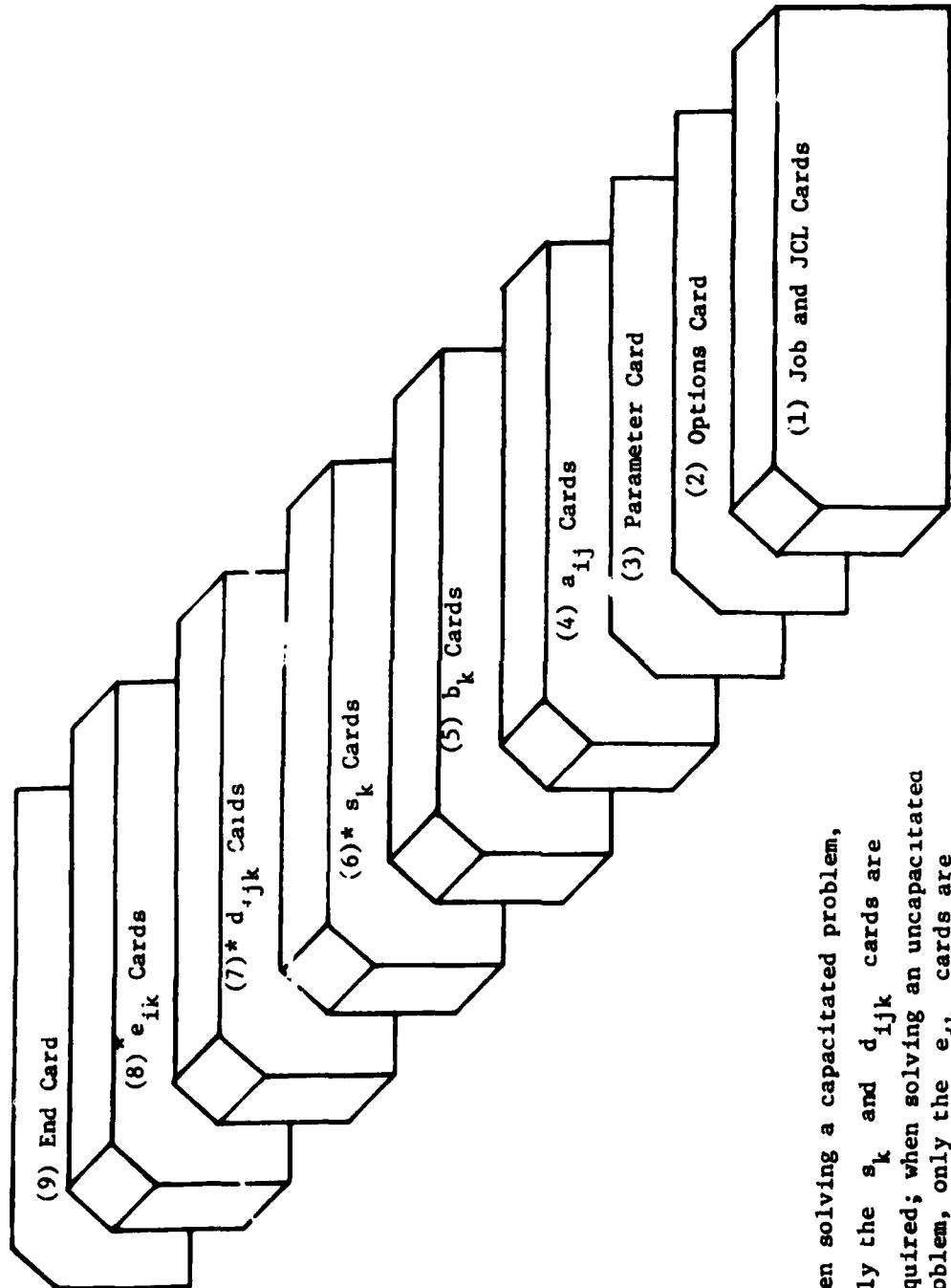
#### 4. User Information

ZIPCAP is designed to solve a multiactivity multifacility capacity-constrained 0-1 assignment problem having formulation (P). Such a problem is referred to as "capacitated" because the optimal solution must satisfy the capacity constraints. However, the program can be used in a situation where the capacity constraints are not applicable, in other words, for an "uncapacitated" problem. The program, in this case, will generate the necessary information to conform to formulation (P).

Figure 3 presents a schematic diagram of the deck structure for using ZIPCAP. The cards include both the job control cards and the input data cards required to solve a problem.

Following is a detailed description of these cards (in the same order as presented in Figure 3); it provides the necessary coding information.

(1) Job and JCL Cards: Figures 4A and 4B show two alternatives for the use of JCL cards.



\*When solving a capacitated problem, only the  $s_k$  and  $d_{ijk}$  cards are required; when solving an uncapacitated problem, only the  $e_{ik}$  cards are required.

Figure 3  
ZIPCAP Data Structure

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## FORTRAN Coding Form

PROGRAM		FORTRAN		JCL	
STATEMENT NO. / NUMBER	TYPE	STATEMENT NO. / NUMBER	TYPE	STATEMENT NO. / NUMBER	TYPE
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80					

## FORTRAN STATEMENT

```

// STANDARD JOBE CARD
// EXEC FORTG2
// FORT.SYSIN DD *
                ↑
(SOURCE PROGRAM)
                ↓
// G0. SYSLIB DD
// DD DSN=GWU. P1IXLIB, DISP=SHR
// G0. SYSIN DD *
                ↓
(DATA CARDS)
                ↓
// 
```

{ JCL CARDS FOR  
TIMEET SUBROUTINE }

Figure 4A

Job and JCL Cards -- Alternative 1

**IBM**

FORTRAN Coding Form

FORTRAN STATEMENT											
3	4	5	6	7	8	9	10	11	12	13	14
/	STANDARD JΦB CARD										
/	EXEC FΦRG3, DSN = 'YΦUR LIBRARY', PRΦG = ZΙΡCAP										
/	FΦRT. SYSSIN DD *										
	↑										
	(SΦURPC1 PΦΦGR( M )										
		↓									

Step 1: Compile and Store Procedure

- 15 -

FORTRAN STATEMENT											
3	4	5	6	7	8	9	10	11	12	13	14
/	STANDARD JΦB CARD										
/	EXEC FΦRG6, DSN = 'YΦUR LIBRARY', PRΦG = ZΙΡCAP										
/	JCL CARDS FOR										
/	TIMET SUBROUTINE										
/	SYSLIB DD										
/	DD DSN = GWU. PLIXLIE, DISP = SHR										
/	SYSSIN DD *										
	↑										
	(DATA CARDS)										
		↓									

Step 2: Load and Go Procedure

Figure 4B

Job and JCL Cards -- Alternative II

Alternative I uses the source program for the compilation and execution procedure. In Alternative II, the program is first compiled and stored, and thereafter, for a given problem, only the execution is necessary.

The program uses the TIMET subroutine from the PLI Library in order to record the elapsed time. This time includes only the execution time of the program and excludes the time needed to read the input and to write the output. The JCL cards needed for this purpose are identified in Figures 4A and 4B. If this subroutine is not available, the corresponding JCL cards are not required. In that case the appropriate source program cards should also be removed. (These cards are identified in the program listing in Appendix A.)

(2) Options Card

The options card is the first input data card and is described in Table 2.

(3) Parameter Card

Table 3 describes the parameter card.

Table 3  
THE PARAMETER CARD

Columns	Format	Name	Definition
1-5	I5	M	Number of designs (m)
6-10	I5	N	Number of activities (n)
11-15	I5	P	Number of facilities(p)

Table 2  
OPTIONS CARD

Columns	Format	Name	Value and Definition	Remarks
i	II	IINPT	1 if input listing desired 0 otherwise	Input listing includes: - parameters m , n , p - $a_{ij}$ , $b_k$ , $s_k$ , $d_{ijk}$ , and $e_{ik}$
2	II	ICAPR	1 if capacity rule to be used 0 otherwise	When solving a capacitated problem (i.e., where capacity constraints are applicable), set ICAPR=1; for an uncapacitated problem, it should be set to 0.
3	II	ISTEP	0 if listing of intermediate steps not desired 1 if summary of intermediate steps desired 2 if detailed intermediate steps desired	<ul style="list-style-type: none"> <li>Summary of intermediate steps includes, for each node, node number, lower bound, upper bound (if applicable), best upper bound, and branching variable. (This information is sufficient to construct a branch-and-bound tree.)</li> <li>Detailed intermediate steps list every computational step at every node. This option is useful only when changing/debugging the program.</li> <li>Even with ISTEP = 0, the output includes the total number of nodes explored.</li> </ul>

Table 2  
(Continued)

Columns	Format	Name	Value and Definition	Remarks
4	I1	IUNCAP	0 if solving a capacitated problem 1 if solving an uncapacitated problem	for IUNCAP=0, ICAPR should be 1, and for IUNCAP=1, ICAPR should be 0. (for IUNCAP=1, if ICAPR=1, the program overrides it and sets ICAPR=0)
5-10	F6.5	EPS	A fractional value if a sub-optimal solution is acceptable. This value should be zero if an optimal solution is desired.	EPS = 0.002 implies that the resulting solution may be suboptimal, but is guaranteed to have a value within 0.2% of the optimal value.
11-20	F10.3	ET	Elapsed time in seconds, if specified, at which the node and bounds information is printed.	<ul style="list-style-type: none"> <li>• Specifying ET is useful in a case where ISTEP = 0 and the program terminates before reaching the final solution.</li> <li>• When applicable, the information printed includes: <ul style="list-style-type: none"> <li>- the node being explored at the specified time</li> <li>- the best upper bound, the corresponding solution, and the node at which this solution was found</li> <li>- detailed steps for the node being explored, by assuming ISTEP=2</li> </ul> </li> <li>• The above information is useful in assessing the extent of the branch-and-bound tree explored until time ET.</li> <li>• This option uses the TIMET subroutine</li> </ul>

(4) to (8) Other Input Data Cards

Table 4 describes the cards for the variable costs  $a_{ij}$ , the fixed costs  $b_k$ , the capacity availabilities  $s_k$ , and the capacity usage values  $d_{ijk}$ . However, if solving an uncapacitated problem, i.e., where capacity constraints are not active, the facility-design values  $e_{ik}$  are required as input data instead of the  $s_k$  and  $d_{ijk}$  values.

The input data values for  $a_{ij}$ ,  $b_k$ ,  $s_k$ , and  $d_{ijk}$  are required in integer format. Any real-valued data can be converted into integer values by scaling, and the optimal solution reconverted to real values later. For example, real-valued data of 215.35, 116.50, ..., 180.61 can be treated as 21535, 11650, ..., 18061; and an optimal value of 584230 obtained from the program would correspond to 5842.30.

(9) End Card

This is the usual end card following input data cards. It contains '//' in the first two columns.

5. Test Problems and Output Illustrations5.1 Test Problem 15.1.1 Problem formulation (stated for ease of reference)

The problem is to find  $x_{ij}$  and  $y_k$  variables that

Table 4  
OTHER INPUT DATA CARDS

Columns	Format	Name	Definition	Remarks
1-10, 11-20, 21-30, etc.	8110	A(I,J)	Variable cost of activity j using design i	The order of the input data a <sub>ij</sub> is a <sub>11</sub> , a <sub>21</sub> , ..., a <sub>m1</sub> ; a <sub>12</sub> , a <sub>22</sub> , ..., a <sub>m2</sub> ; ..., a <sub>1n</sub> , a <sub>2n</sub> , ..., a <sub>mn</sub>
1-10, 11-20, 21-30, etc.	8110	B(K)	Fixed cost of facility k	The order for the b <sub>k</sub> is b <sub>1</sub> , b <sub>2</sub> , ..., b <sub>p</sub>
1-10, 11-20, 21-30, etc.	8110	S(K)	Capacity available at facility k	. The order for the s <sub>k</sub> is s <sub>1</sub> , s <sub>2</sub> , ..., s <sub>p</sub> . This input is required when solving a capacitated problem only, i.e., for IUNCAP=0. For an uncapacitated problem, i.e., for IUNCAP=1, this information is generated by the program.

Table 4  
(Continued)

Columns	Format	Name	Definition	Remarks
1-10, 11-20, 21-30, etc.	8I10	D(I,J,K)	Capacity required at facility $k$ for activity $j$ uses design $i$ . A $d_{ijk}$ value is positive if design $i$ uses facility $k$ ; otherwise it is 0.	<ul style="list-style-type: none"> <li>The order of input data is all <math>d_{ijk}</math> for <math>k=1</math>, then all <math>d_{ijk}</math> for <math>k=2</math>, etc., starting on a new card when the value of <math>k</math> changes.</li> <li>The order of input for a facility <math>k</math> is <math>d_{11k}, d_{21k}, \dots, d_{m1k};</math>  <math>d_{12k}, d_{22k}, \dots, d_{m2k};</math>  <math>\dots</math>  <math>d_{1nk}, d_{2nk}, \dots, d_{mnk}.</math></li> <li>This input, similar to <math>s_k</math>, is required for capacitated problems only, i.e., for IUNCAP=0. This information is generated by the program for an uncapacitated problem, i.e., for IUNCAP=1.</li> </ul>

Table 4  
(Continued)

Columns	Format	Name	Definition	Remarks
1,2,3,...	80I1	E(I,K)	$e_{ik} = 1$ if design $i$ uses facility $k$ $= 0$ otherwise	<ul style="list-style-type: none"> <li>The order of input for the <math>e_{ik}</math> is <math>e_{11}, e_{21}, \dots, e_{m1};</math></li> <li><math>e_{12}, e_{22}, \dots, e_{m2}; \dots,</math></li> <li><math>e_{1p}, e_{2p}, \dots, e_{mp}.</math></li> </ul> <p>This input is required when solving an uncapacitated problem only, i.e., for IUNCAP=1. The program, then generates <math>d_{ijk}</math> and <math>s_k</math> as follows:</p>

$$\begin{aligned}
 d_{ijk} &= 1 \quad \text{for all } j \text{ if} \\
 e_{ik} &= 1 \\
 &= 0 \quad \text{for all } j \text{ if} \\
 e_{ik} &= 0
 \end{aligned}$$

$$s_k = n \quad \text{for all } j$$

- $e_{ik}$  is generated by the program for a capacitated problem, as follows:

$$\begin{aligned}
 e_{ik} &= 1 \quad \text{if } \sum_j d_{ijk} \text{ is} \\
 &\quad \text{positive } (>0) \\
 &= 0 \quad \text{otherwise}
 \end{aligned}$$

$$(P) \left\{ \begin{array}{l} \text{Minimize} \quad \sum_{j=1}^n \sum_{i=1}^m a_{ij} x_{ij} + \sum_{k=1}^p b_k y_k \quad (1) \\ \text{subject to} \quad \sum_{i=1}^m x_{ij} = 1 \quad j=1, \dots, n \quad (2) \\ \quad \sum_{j=1}^n \sum_{i=1}^m d_{ijk} x_{ij} \leq s_k y_k \quad k=1, \dots, p \quad (3) \\ \quad x_{ij}, y_k = 0 \text{ or } 1, \end{array} \right.$$

where  $m = 3$ ,  $n = 4$ ,  $p = 5$ .

The  $a_{ij}$  values are:

		j			
		1	2	3	4
i	1	9,847	15,718	4,450	8,925
	2	10,192	10,542	4,422	8,118
	3	11,019	9,750	4,609	8,337

$$b_1 = 1750 \quad b_2 = 2000 \quad b_3 = 1750 \quad b_4 = 1350 \quad b_5 = 1000$$

$$s_1 = 400 \quad s_2 = 400 \quad s_3 = 1000 \quad s_4 = 400 \quad s_5 = 400$$

The  $d_{ijk}$  values are as follows:

		j			
		1	2	3	4
i	k=1	0	0	0	0
	2	186	60	261	148
	3	154	38	175	100

		j			
		1	2	3	4
k=2	1	0	0	0	0
	2	0	0	0	0
	3	154	38	175	100

		j			
		1	2	3	4
k=3	1	525	144	1,011	396
	2	186	60	261	148
	3	154	38	175	100

		j			
		1	2	3	4
k=4	1	0	0	0	0
	2	0	0	0	0
	3	154	38	175	100

		j			
		1	2	3	4
k=5	1	0	0	0	0
	2	186	60	261	148
	3	154	38	175	100

5.1.2 Coded Input: Figure 5 shows the coded information for test problem 1. Various segments in this figure are numbered in parentheses to correspond to the deck structure of Figure 3.

IBM

## FORTRAN Coding Form

GX287327-6 U/M 050-  
Printed in U.S.A.

STATEMENT NUMBER	FORTRAN STATEMENT		P. ADDRESS DISPLAYED	JOURNAL	FUNCTION	CARD ELECTRIC NUMBER
	1	2				
11	STANDARD J48 CARD					
11	EXEC F4PG6,DSN='YFOR LIBRARY',PRG=21PCAP					
11	GΦ·SYSLIB DD					
11	DD DSN=GWWU·PLIXLIB,DIS P=SHR					
11	GΦ·SYSIN DD *					
111	3 4 5					
111	9847 10192 11019	15718 10542	9750	4450	4422	(4)
111	4609 8925	8118 8337				(5)
111	1750 2000	1750 1350	1000			(6)
111	400 400	1000 400	400	38	0	(7)
111	0 196	154 0	60			
111	175 0	148 100	0			
111	0 0	154 0	0	38	0	
111	175 0	0 100	0			
111	525 186	154 144	60	38	1011 264	
111	175 396	148 100				
111	0 0	154 0	0	38	0	
111	175 0	0 100	0			
111	0 186	154 0	60	38	0	
111	175 0	148 100	0			

Figure 5  
Coded Input Data for Test Problem 1

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5.1.3 Annotated Output: The computer output for this test problem is presented in Appendix C. Most segments of this output are self-explanatory. The order of these segments is lettered in circles as follows:

- (a) Options selected -- information given on options card
- (b) m, n, and p -- information from parameter card. This segment is printed if the option IINPT=1
- (c) Other input data, i.e.,  $a_{ij}$ ,  $b_k$ ,  $s_k$  and  $d_{ijk}$  for a capacitated problem; and  $e_{ik}$  as generated by the program. For an uncapacitated problem, input data  $a_{ij}$ ,  $b_k$ , and  $e_{ik}$ ; and  $s_k$  and  $d_{ijk}$  as generated by the program. This segment is printed if the option IINPT=1
- (d) Intermediate steps' display, depending on the value of the option ISTEP. No display for ISTEP=0, summary for ISTEP=1, and detailed steps for ISTEP=2
- (e) Elapsed time in seconds to execute the program, excluding the time to read the input data and to write the input and output. It includes the time to print intermediate steps if option ISTEP=1 or 2
- (f) Total number of nodes explored for the problem solved
- (g) Optimal solution showing the  $x_{ij}$  variables with value 1,  $y_k$  variables with value 1 or 0, and the value of the objective function. If the problem does not have a feasible solution, this segment prints "Problem does not have a feasible solution because the capacity constraints cannot be satisfied."

The node and bounds information at a specified time is displayed only if option RT is assigned a particular value.

### 5.2 Test Problem 2

This is a capacitated problem having 5 designs, 4 activities and 8 facilities.

The problem formulation and coded input are similar to that of test problem 1.

The annotated output (which includes the listing of input data) is presented in Appendix D. The annotations are the same as those already described for test problem 1. Note that the options selected include ISTEP=0, hence no intermediate steps; and EPS=0.002, implying that the solution may be suboptimal within + 0.2% .

### 5.3 Test Problem 3

Although ZIPCAP is designed for capacitated problems, it may be used to solve uncapacitated problems as a special case. For demonstration, this test problem is uncapacitated, having 10 designs, 8 activities, and 8 facilities. The input data include the options card, the parameter card, input data values for  $a_{ij}$ ,  $b_k$ , and  $e_{ik}$ . The program generates the  $d_{ijk}$  and  $s_k$  values.

Appendix E shows the annotated output for this problem. Note that the options selected include IINPT=1, hence the data listing; ISTEP=0, therefore no intermediate steps; ICAPR=0 and IUNCAP=1 being an uncapacitated problem. Note that option ET has been specified a value of 30.0 . Thus, if the optimal solution was not reached within 30 seconds, the node and bounds information at time ET would have been displayed.

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**APPENDIX A**  
**ZIPCAP LISTING**

FORTRAN IV G LEVEL 21

MAIN

DATE = 80206

14/35/07

```

C      ZIPCAP, A ZERO-ONE INTEGER PROGRAM IS DESIGNED      00000010
C      TO SOLVE MULTIACTIVITY MULTIFACILITY CAPACITY-      00000015
C      CONSTRAINED PROBLEMS HAVING VARIABLE AND FIXED      00000020
C      COSTS. IT ALSO SOLVES UNCAPACITATED PROBLEMS AS A  00000030
C      SPECIAL CASE                                         00000040
0001    INTEGER    D(35,35,30), A(35,35), CX(35,35), E(35,30), 00000050
        1       B(30), BSOLX(35), BSOLY(30), FLB(30), FIX(35), FIXI(35), 00000060
        2       FUB(30), S(30), SOLX(35), STX(1225)                00000070
0002    REAL       MINC(35), NMINC(35)                         00000080
0003    DIMENSION C(35,35), DIFBR(35), KT2(35), MIND(35)        00000090
0004    INTEGER    BRO, BR1, FC, FCUB, P                         00000110
0005    REAL       LOWB, MAXDIF, MINSC                         00000120
C      *****OPTIONS AVAILABLE: IINPT, ICAPP, ISTEP, IUNCAP, EPS 00000130
C      IINPT=1 IF INPUT LISTING DESIRED; 0 OTHERWISE          00000140
C      ICAPR=1 IF CAPACITY RULE TO BE USED; 0 OTHERWISE       00000150
C      ISTEP=0 IF LISTING OF INTERMEDIATE STEPS             00000160
C      NOT DESIRED. ISTEP=1 IF SUMMARY OF BRANCH &           00000170
C      BOUND NODES DESIRED. ISTEP=2 IF DETAILED             00000180
C      LISTING OF INTERMEDIATE STEPS DESIRED.                 00000190
C      IUNCAP=1 IF SOLVING AN UNCAPACITATED PROBLEM,         00000200
C      0 OTHERWISE.                                         00000210
C      EPS= A FRACTIONAL VALUE IF SUBOPTIMAL                 00000220
C      SOLUTION DESIRED, E.G., EPSILON AS 0.005              00000230
C      IMPLIES SOLUTION TO BE WITHIN +0.5 PERCENT           00000240
C      OF THE OPTIMAL SOLUTION. EPS=0.0 IF OPTIMAL            00000250
C      SOLUTION DESIRED.                                     00000253
C      ET= ELAPSED TIME IN SECONDS, IF SPECIFIED, AT          00000256
C      WHICH THE NODE AND BOUNDS RELATED INFORMATION          00000260
C      IS PRINTED. THIS IS USEFUL IN A SITUATION IF          00000263
C      ISTEP=0 AND THE PROGRAM TERMINATES BEFORE             00000266
C      REACHING THE FINAL SOLUTION.                          00000270
C      *****READ INPUT DATA*****                           00000273
0006    READ 10, IINPT, ICAPR, ISTEP, IUNCAP, EPS, ET          00000280
0007    10 FORMAT (4I1, F6.5, F10.3)                         00000290
C      M= NUMBER OF DESIGNS                            00000300
C      N= NUMBER OF ACTIVITIES                         00000310
C      P= NUMBER OF FACILITIES                        00000320
0008    RE>D 20,M,N,P                                     00000330
0009    20 FORMAT (3I5)                                    00000340
C      A(I,J): VARIABLE COST MATRIX                  00000350
0010    READ 30, ((A(I,J), I=1,M),J=1,N)               00000360
0011    30 FORMAT (IB10)                                00000370
C      B(K): FIXED COST VECTOR                      00000380
0012    READ 30, (BIK),K=1,P                           00000390
0013    IF (IUNCAP.EQ.1) GO TO 40                     00000400
C      S(IK): CAPACITY LIMIT VECTOR; REQUIRED ONLY     00000410
C      IF IUNCAP=0                                     00000420
0014    READ 30, (ISIK),K=1,P                           00000430
C      D(I,J,K): CAPACITY USAGE MATRIX; REQUIRED      00000440
C      ONLY IF IUNCAP=0                               00000450
0015    DO 32 K=1,P                                     00000460
0016    READ 30,((D(I,J,K), I=1,M),J=1,N)            00000470
0017    32 CONTINUE                                 00000480
0018    DO 37 K=1,P                                     00000490
0019    DO 37 I=1,M                                    00000500
0020    IF (D(I,1,K).EQ.0) GO TO 35                  00000510
0021    E(I,K)=1                                    00000520
0022    GO TO 37                                    00000530

```

FORTRAN IV G LEVEL 21	MAIN	DATE = 80206	14/35/07
0023	35 E(I,J)=0		00000540
0024	37 CONTINUE		00000550
0025	GO TO 90		00000560
C	E(I,K): DESIGN-FACILITY MATRIX; REQUIRED ONLY		00000570
C	IF IUNCAP=1		00000580
0026	40 READ 45,((E(I,K),I=1,M),K=1,P)		00000590
0027	45 FORMAT (180I1)		00000600
0028	DO 80 K=1,P		00000610
0029	S(K)=N		00000620
0030	DO 75 I=1,M		00000630
0031	IF (E(I,K).EQ.1) GO TO 65		00000640
0032	DO 60 J=1,N		00000650
0033	D(I,J,K)=0		00000660
0034	60 CONTINUE		00000670
0035	GO TO 75		00000680
0036	65 DO 70 J=1,N		00000690
0037	D(I,J,K)=1		00000700
0038	70 CONTINUE		00000710
0039	75 CONTINUE		00000720
0040	80 CONTINUE		00000730
C	*****PRINT INPUT DATA*****		00000740
0041	90 PRINT 95, IINPT, ICAPR, ISTEP, IUNCAP, EPS, ET		00000750
0042	95 FORMAT (1'1',' OPTIONS SELECTED : IINPT=',I1,		00000760
	1 ' ICAPR=',I1, ' ISTEP=',I1, ' IUNCAP=',I1,		00000770
	2 ' EPS=',F8.5, ' ET=', F10.3//)		00000780
0043	IF (IINPT.EQ.0) GO TO 168		00000790
0044	PRINT 100,M,N,P		00000800
0045	100 FORMAT (1'0', T55, 'INPUT DATA',/1X, T55, '-----',//1X,		00000810
	1T41, 'NUMBER OF DESIGNS (M)=', 4X,I4//1X,T41,		00000820
	2'NUMBER OF ACTIVITIES (N)=', 1X,I4//1X, T41,		00000830
	3'NUMBER OF FACILITIES (P)=',1X, I4//)		00000840
0046	PRINT 105		00000850
0047	105 FORMAT (1 4X, 'VARIABLE COST MATRIX A(I,J)',/4X,		00000860
	1'-----',/)		00000870
0048	DO 110 I=1,M		00000880
0049	110 PRINT 115, I, (A(I,J),J=1,N)		00000890
0050	115 FORMAT (1'0', T6, 'I=', I3, 4X,8I13, 4(/, 14X,8I13))		00000900
0051	PRINT 120		00000910
0052	120 FORMAT(1'0',//4X,'FIXED COST VECTOR B(K)',/4X,		00000920
	1'-----',/)		00000930
0053	PRINT 122, (B(K),K=1,P)		00000940
0054	122 FORMAT (1'0', T15, 8I13, 3(/, 14X,8I13))		00000950
0055	PRINT 125		00000960
0056	125 FORMAT(1'0',//4X,'CAPACITY LIMIT VECTOR S(K)',/4X,		00000970
	1'-----',/)		00000980
0057	PRINT 128, (S(K),K=1,P)		00000990
0058	128 FORMAT (1'0', T15, 8I13, 3(/, 14X,8I13))		00001000
0059	PRINT 130		00001010
0060	130 FORMAT(1'0',//4X,'CAPACITY USAGE MATRIX D(I,J,K)',/4X,		00001020
	1'-----',/)		00001030
0061	DO 150 K=1,P		00001040
0062	PRINT 135,K		00001050
0063	135 FORMAT (1'0',//5X,'K=',I3//)		00001060
0064	DO 145 I=1,M		00001070
0065	PRINT 140,I,(D(I,J,K), J=1,N)		00001080
0066	140 FORMAT (1'0', T6, 'I=', I3, 4X,8I13, 4(/, 14X,8I13))		00001090
0067	145 CONTINUE		00001100
0068	150 CONTINUE		00001110

FORTRAN IV G LEVEL 21

	MAIN	DATE = 80206	14/35/07
0069	PRINT 155		00001120
0070	155 FORMAT('0',//4X,'DESIGN-FACILITY MATRIX E(I,K)',/4X, 1'-----',/)		00001130
0071	DO 160 I=1,M		00001140
0072	PRINT 158, I, IE(I,K),K=1,P)		00001150
0073	158 FORMAT ('0', T6, 'I=', I3, 4X,8I13, 3(/, 14X,8I13))		00001160
0074	160 CONTINUE		00001180
0075	168 IF (ISTEP.EQ.0) GO TO 190		00001190
0076	IF (ISTEP.EQ.1) GO TO 175		00001200
0077	PRINT 170		00001210
0078	170 FORMAT ('0',//55X,'DETAILED LISTING OF STEPS',/)		00001220
0079	GO TO 190		00001230
0080	175 PRINT 180		00001240
0081	180 FORMAT ('0',//55X,'SUMMARY OF STEPS',/)		00001250
	C *****INITIALIZE*****		00001260
0082	190 BUB=999999.		00001270
0083	BUBS= BUB/ (1.0+EPS)		00001280
0084	NSX=0		00001290
0085	NOD=1		00001310
0086	IBMOD=1		00001315
0087	INFT=0		00001320*
0088	INSET=0		00001330*
0089	DO 205 J=1,N		00001390
0090	FIX(J)=0		00001400
0091	KT2(J)=0		00001410
0092	DO 205 I=1,M		00001420
0093	CX(I,J)=0		00001430
0094	205 CONTINUE		00001433
0095	LQ1=0		00001436
0096	LQ2=0		00001440
0097	LR2=0		00001443
0098	CALL TIMET(1TO)		00001445*
0099	IF(NSX.EQ.0) GO TO 283		00001447
	C CX(I,J) CONTAINS FIXED AND FREE X(I,J) VARIABLES.		00001450
	C STX(INS) CONTAINS FIXED X(I,J) VARIABLES.		00001460
	C CX(I,J) AND STX(INS) ARE UPDATED BY THE CAPACITY		00001460
	C RULE, THE BOUNDING RULE, AND THE RULE FOR		00001490
	C BRANCHING AND BACKTRACKING.		00001500
	C IN CX(I,J) A FIXED VARIABLE IS RECORDED AS 1 OR		00001505
	C 2, AND A FREE VARIABLE AS 0.		00001510
	C A VALUE OF 1 IMPLIES THAT THAT PARTICULAR VARIABLE		00001515
	C IS FIXED, AND FIX(J) IS SET EQUAL TO 1 IMPLYING		00001520
	C THAT COLUMN J HAS A FIXED VARIABLE OF VALUE 1.		00001525
	C A VALUE OF 2 IMPLIES THAT THAT PARTICULAR VARIABLE		00001530
	C SHOULD NOT BE CONSIDERED FOR CURRENT COMPUTATIONS.		00001535
	C AN X(I,J) RECORDED IN CX(I,J) AS 1 DUE TO THE		00001540
	C BRANCHING RULE IS RECORDED IN STX(INS) AS X*100+J.		00001545
	C AN X(I,J) RECORDED IN CX(I,J) AS 1 DUE TO THE		00001550
	C CAPACITY RULE OR THE BOUNDING RULE IS RECORDED IN		00001555
	C STX(INS) AS (X*100+J)+1000000.		00001560
	C AN X(I,J) RECORDED IN CX(I,J) AS 2 IS RECORDED IN		00001565
	C STX(INS) AS -(X*100+J)-1000000.		00001570
0100	210 IF (ISTEP.EQ.0) GO TO 225		00001580
0101	215 PRINT 220,NUD		00001590
0102	220 FORMAT ('0',// 6X,'NODE NUMBER',I5/)		00001600
	C *****UPDATE CX(I,J) FOR BRO*****		00001610
	C BRO IS THE RIGHT BRANCHING VARIABLE		00001615
0103	225 LX=BRO		00001620

\*Relevant to the TIMET Subroutine

FORTRAN IV G LEVEL 21	MAIN	DATE = 80206	14/35/07
0104	IX=LX/100		00001630
0105	JX=LX-IX*100		00001640
0106	CX(IX,JX)=2		00001650
0107	KT2(JX)=KT2(JX)+1		00001660
0108	FIX(JX)=0		00001720
0109	LQ1=LQ1-1		00001725
0110	IF (KT2(JX).LT.(M-1)) GO TO 270		00001730
0111	DO 255 I=1,M		00001740
0112	IF (CX(I,JX).EQ.2) GO TO 255		00001750
0113	CX(I,JX)=1		00001760
0114	NSX=NSX+1		00001763
0115	STX(NSX)= (I*100+JX)+1000000		00001766
0116	FIX(JX)=1		00001770
0117	LQ1=LQ1+1		00001780
0118	FIXI(JX)=I		00001790
0119	GO TO 270		00001800
0120	255 CONTINUE		00001810
0121	270 LQ2=0		00001820
0122	LR2=0		00001825
0123	GO TO 283		00001830
0124	272 IF (ISTEP.EQ.0) GO TO 276		00001840
0125	PRINT 220,NOD		00001850
	C ***** UPDATE CX(I,J) FOR BR1***** C BR1 IS THE LEFT BRANCHING VARIABLE		00001853
0126	276 LQ2=0		00001860
0127	LR2=0		00001866
0128	LX=BR1		00001870
0129	IX=LX/100		00001875
0130	JX=LX-IX*100		00001880
0131	CX(IX,JX)=1		00001885
0132	FIX(JX)=1		00001890
0133	LQ1=LQ1+1		00001892
0134	DO 279 I=1,M		00001895
0135	IF (IX.EQ.I) GO TO 281		00001897
0136	279 CONTINUE		00001900
0137	281 FIXI(JX)=IX		00001902
0138	283 IF (ISTEP.NE.2) GO TO 300		00001905
0139	285 DO 295 I=1,M		00001910
0140	PRINT 290, I,(CX(I,J),J=1,N)		00001920
0141	FORMAT (/5X,'CX(I,J)',4X,'I=',I3,2X, 20I4/23X, 20I4)		00001930
0142	295 CONTINUE		00001940
0143	PRINT 297,(FIX(J),J=1,N)		00001950
0144	FORMAT (/5X,'FIX(J)',12X, 20I4/23X, 20I4)		00001960
	C ***** APPLY CAPACITY RULE***** C AND UPDATE CX(I,J) AND STX(INS).		00001970
0145	300 IF (IUNCAP.EQ.1) GO TO 2300		00002010
0146	305 IF (ICAPR.EQ.0) GO TO 2300		00002020
0147	310 DO 2000 K=1,P		00002030
	C FIND THE SUM OF MINIMUM D(I,J,K) OVER EACH J FOR A C GIVEN K, I.E., MIND= SUM OF MIND(J)		00002040
0148	MIND=0		00002050
0149	DO 900 J=1,N		00002070
0150	IF(FIX(J).EQ.0) GO TO 350		00002080
	IF FIX(J)=1, SET MIND(J)=D(I,J,K) FOR CX(I,J)=1 AND MOVE TO NEXT COLUMN J		00002090
0151	INDI=FIXI(J)		00002110
0152	MIND(J)=D(INDI,J,K)		00002120
0153	GO TO 800		00002130

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0154	350	LK=0		00002160
0155		I=1		00002170
0156		MIND(J)=D(I,J,K)		00002180
	C	SKIP D(I,J,K) WHEN CX(I,J)=2 & MOVE TO NEXT ROW I		00002190
0157	400	IF(CX(I,J).EQ.2) GO TO 600		00002200
0158	500	IF(D(I,J,K).LT.MIND(J)) MIND(J)=D(I,J,K)		00002210
0159		GO TO 700		00002220
0160	600	LK=LK+1		00002230
0161		IF(I.GT.LK) GO TO 700		00002240
0162		I=I+1		00002250
0163		MIND(J)=D(I,J,K)		00002260
0164		GO TO 750		00002270
0165	700	I=I+1		00002280
0166	750	IF(I.LE.M) GO TO 400		00002290
0167	800	MINSD=MINSD+MIND(J)		00002300
0168	900	CONTINUE		00002310
0169	910	IF (ISTEP.NE.2) GO TO 980		00002320
0170		PRINT 950, K, MINSD,(MIND(J),J=1,N)		00002330
0171	950	FORMAT('0',*K, MINSD,(MIND(J),J=1,N)*,10I10)		00002340
	C	FIND BALANCE AVAILABLE CAPACITY IBALD FOR A GIVEN K		00002350
	C	IF IBALD IS NEGATIVE, THEN BACKTRACK.		00002360
0172	980	IBALD=S(K)-MINSD		00002380
0173	1000	IF (IBALD.LT.0) GO TO 6200		00002390
0174		DO 1500 J=1,N		00002400
	C	SKIP COLUMN J IF FIX(J)=1		00002410
0175		IF (FIX(J).EQ.1) GO TO 1500		00002420
0176		DO 1300 I=1,M		00002430
	C	SKIP ROW I IF CX(I,J)=2		00002440
0177	1100	IF(CX(I,J).EQ.2) GO TO 1300		00002450
	C	COMPUTE DIFFERENCE BETWEEN D(I,J,K) AND MIND(J).		00002470
	C	IF IT IS MORE THAN AVAILABE BALANCE, SET CX(I,J)=2		00002480
0178	1200	IDIFD=D(I,J,K)-MIND(J)		00002490
0179		IF ((IDIFD-IBALD).LE.0) GO TO 1300		00002510
0180		CX(I,J)=2		00002520
0181		NSX=NSX+1		00002523
0182		STX(NSX)=-(I*100+J)-1000000		00002526
	C	LQ2 COUNTS THE NUMBER OF CX(I,J) VALUES SET EQUAL		00002530
	C	TO 2 IN A CYCLE		00002540
0183		LQ2=LQ2+1		00002550
	C	KT2(J) KEEPS AN ACCOUNT OF CX(I,J) VALUES SET EQUAL		00002560
	C	TO 2 FOR COLUMN J		00002570
0184		KT2(J)=KT2(J)+1		00002580
	C	FOR COLUMN J, IF ALL BUT ONE CX(I,J) VALUES ARE		00002590
	C	EQUAL TO 2, SET THAT CX(I,J)=1 & SET FIX(J)=1		00002600
0185		IF(KT2(J).LT.(M-1)) GO TO 1300		00002610
0186	DO 1250	LR=1,M		00002620
0187		IF(CX(LR,J).EQ.2) GO TO 1250		00002630
0188		CX(LR,J)=1		00002640
0189		NSX=NSX+1		00002643
0190		STX(NSX)=(LR*100+J)+1000000		00002646
0191		FIX(J)=1		00002650
	C	LQ1 KEEPS AN ACCOUNT OF COLUMNS FOR WHICH FIX(J)=1		00002655
0192		LQ1=LQ1+1		00002660
	C	FIXI(J) SPECIFIES INDEX I FOR WHICH FIX(J)=1		00002662
0193		FIXI(J)=LR		00002665
0194		GO TO 1500		00002670
0195	1250	CONTINUE		00002680
0196	1300	CONTINUE		00002690

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0197      1500 CONTINUE          00002700
0198      1800 IF (ISTEP.NE.2) GO TO 2000 00002710
0199      PRINT 1900, K, LQ2, LQ1 00002720
0200      1900 FORMAT ('0', 'K=0,I3,1 Q2=0,I3,1 LQ1=0, I3) 00002730
0201      DO 1930 I=1,M 00002740
0202      PRINT 290, I,(CX(I,J),J=1,N) 00002750
0203      1930 CONTINUE 00002770
0204      PRINT 297, (IFIX(J),J=1,N) 00002780
0205      2000 CONTINUE 00002800
C           A CYCLE EXAMINES ALL THE FACILITIES. 00002803
C           IF IN A CYCLE, THE CAPACITY RULE RESULTS IN SETTING
C           ADDITIONAL CX(I,J) VALUES EQUAL TO 2, THEN REPEAT
C           THE CYCLE. BUT IF FIX(J)=1 FOR ALL J, THEN DO NOT
C           REPEAT THE CYCLE. 00002810
0206      IF (LQ1.EQ.N) GO TO 2300 00002840
0207      IF (LU2.EQ.LR2) GO TO 2300 00002845
0208      2200 LR2=LQ2 00002860
0209      GO TO 300 00002870
C           ***** SOLVE (LAGRANGIAN) RELAXED PROBLEM*****
C           FIND FLB(K) — VECTOR OF FACILITIES FOR COMPUTING
C           C(I,J) MATRIX & LOWER BOUND. IT HAS VALUE 1 IF A
C           FACILITY IS USED, OTHERWISE IT HAS 0 VALUE. 00002880
0210      2300 DO 2500 K=1,P 00002910
0211      FLB(K)=0 00002930
0212      2500 CONTINUE 00002940
0213      DO 3000 J=1,N 00002950
0214      IF (FIX(J).EQ.0) GO TO 3000 00002960
0215      INDI=FIXI(J) 00002970
0216      DO 2550 K=1,P 00002990
0217      IF (E(INDI,K).EQ.0) GO TO 2550 00003000
0218      IF (FLB(K).EQ.1) GO TO 2550 00003010
0219      FLB(K)=1 00003020
0220      2550 CONTINUE 00003030
0221      3000 CONTINUE 00003060
0222      IF (ISTEP.NE.2) GO TO 3150 00003070
0223      PRINT 3100, (FLB(K),K=1,P) 00003080
0224      3100 FORMAT('0',(FLB(K),K=1,P) ", 20I4/16X,20I4) 00003090
C           COMPUTE COST MATRIX C(I,J) FOR THE RELAXED PROBLEM 00003100
0225      3150 DO 3400 J=1,N 00003110
0226      DO 3300 I=1,M 00003120
0227      BSUM=0.0 00003130
0228      DO 3200 K=1,P 00003140
0229      IF (FLB(K).EQ.1) GO TO 3200 00003150
0230      IF (E(I,K).EQ.0) GO TO 3200 00003160
0231      BSUM=BSUM+(B(K) * (FLOAT(D(I,J,K))/ FLOAT(S(K)))) 00003170
0232      3200 CONTINUE 00003180
0233      3250 C(I,J)=A(I,J)+BSUM 00003190
0234      3300 CONTINUE 00003200
0235      3400 CONTINUE 00003210
0236      IF (ISTEP.NE.2) GO TO 3445 00003220
0237      DO 3430 I=1,M 00003230
0238      PRINT 3420, I, (C(I,J),J=1,N) 00003250
0239      3420 FORMAT (1/5X, 'C(I,J)', 5X, 'I=0, 13,2X, 5F15.4,
1           6(1/23X, 5F15.4)) 00003260
0240      3430 CONTINUE 00003270
C           FIND SUM OF MINIMUM C(I,J) VALUES OVER EACH J,
C           I.E., MINSC=SUM OF MINC(J). 00003290
C           IF FIX(J)=1, THEN MINC(J)=C(I,J) WHERE CX(I,J)=1 00003300
C           IF FIX(J)=1, THEN MINC(J)=C(I,J) WHERE CX(I,J)=1 00003310

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	MAIN	DATE = 80206	14/35/07
0241	3445 MINSC=0.0		00003320
0242	DO 3900 J=1,N		00003340
0243	IF (F1X(J).EQ.0) GO TO 3500		00003350
0244	INDI=FIXI(J)		00003360
0245	MINC(J)=C(INDI,J)		00003370
0246	SOLX(J)=INDI		00003380
0247	GO TO 3850		00003410
0248	3500 LK=0		00003430
0249	I=1		00003440
C	SKIP C(I,J) ELEMENT IF CX(I,J)=2 & MOVE TO NEXT I		00003470
0250	3550 IF (CX(I,J).EQ.2) GO TO 3700		00003480
0251	IF ((I-LK).EQ.1) GO TO 3600		00003485
0252	IF (C(I,J).GE.MINC(J)) GO TO 3750		00003490
0253	3600 MINC(J)=C(I,J)		00003500
0254	IMIN=I		00003510
0255	GO TO 3750		00003520
0256	3700 LK=LK+1		00003530
0257	3750 I=I+1		00003590
0258	3800 IF (I.LE.M) GO TO 3550		00003600
0259	SOLX(J)=IMIN		00003610
0260	3850 MINSC=MINSC+MINC(J)		00003620
0261	3900 CONTINUE		00003630
0262	IF (ISTEP.NE.2) GO TO 3940		00003640
0263	DO 3920 J=1,N		00003650
0264	PRINT 3910, J,MINC(J),SOLX(J)		00003660
0265	3910 FORMAT ('0', 'J,MINC(J),SOLX(J)', 15,F15.4,I6)		00003670
0266	3920 CONTINUE		00003680
C	COMPUTE FIXED COST FC FOR LOWER BOUND		00003710
0267	3940 FC=0		00003720
0268	DO 4000 K=1,P		00003730
0269	IF (FLB(K).EQ.0) GO TO 4000		00003740
0270	FC=FC+B(K)		00003750
0271	4000 CONTINUE		00003760
C	***** COMPUTE LOWER BOUND LOWB *****		00003770
0272	4050 LOWB=MINSC+FC		00003780
0273	IF (ISTEP.EQ.0) GO TO 4150		00003790
0274	PRINT 4120, MINSC, FC, LOWB		00003800
0275	4120 FORMAT ('0', ' MINSC, FC, LOWB ', F15.4, I15, F15.4)		00003810
C	COMPARE LOWER BOUND WITH BEST UPPER BOUND STAR		00003820
C	BUBS WHICH EQUALS BUB/(1+EPS). IF LOWB IS		00003830
C	GREATER THAN OR EQUAL TO BUBS, THEN BACKTRACK		00003840
0276	4150 IF (LOWB.GE.BUBS)GO TO 6200		00003850
C	CHECK IF CURRENT SOLUTION SATISFIES CAPACITY		00003880
C	CONSTRAINTS		00003890
0277	~200 IF (IUNCAP.EQ.1) GO TO 4420		00003900
0278	4210 DO 4400 K=1,P		00003910
0279	NSUMD=0		00003920
0280	DO 4300 J=1,N		00003930
0281	IX=SOLX(J)		00003950
0282	NSUMD=NSUMD+D(IX,J, K)		00003960
0283	4300 CONTINUE		00003970
0284	IF (ISTEP.NE.2) GO TO 4320		00003980
0285	PRINT 4310, K,NSUMD		00003990
0286	4310 FORMAT ('0', 'K,NSUMD', 2I10)		00004000
0287	4320 IF(NSUMD.LE.SIK)) GO TO 4400		00004010
0288	GO TO 5100		00004020
0289	4400 CONTINUE		00004030
C	***** COMPUTE UPPER BOUND UPB IF CAPACITY CONSTRAINTS		00004040

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C                  ARE SATISFIED.                  00004050
C                  UPB=SUM OF A(I,J)+FIXED COST FCUB BASED ON                  00004060
C                  SOLUTION VECTOR SOLX(J)                  00004070
C                  VECTOR OF FACILITIES FOR UPPER BOUND FUB(K) HAS                  00004080
C                  VALUES 1 OR 0 BASED ON FACILITY USED OR OTHERWISE                  00004090
C
0290                  4420 DO 4450 K=1,P                  00004100
0291                  FUB(K)=0                  00004110
0292                  4450 CONTINUE                  00004120
0293                  NSUMA=0                  00004130
0294                  FCUB=0                  00004140
0295                  4500 DO 4650 J=1,N                  00004150
0296                  IX=SOLX(J)                  00004170
0297                  NSUMA=NSUMA+A(IX,J)                  00004180
0298                  4550 DO 4600 K=1,P                  00004190
0299                  IF(E(IX,K).EQ.0) GO TO 4600                  00004200
0300                  IF(FUB(K).EQ.1) GO TO 4600                  00004210
0301                  FUB(K)=1                  00004220
0302                  FCUB=FCUB+B(K)                  00004230
0303                  4600 CONTINUE                  00004240
0304                  4650 CONTINUE                  00004250
0305                  IF (ISTEP.NE.2) GO TO 4700                  00004260
0306                  PRINT 4660, (FUB(K),K=1,P)                  00004270
0307                  4660 FORMAT("0", "(FUB(K),K=1,P)", " 20I4/16X,20I4)                  00004280
0308                  4700 UPB=NSUMA+FCUB                  00004290
0309                  4708 IF (ISTEP.EQ.0) GO TO 4750                  00004300
0310                  PRINT 4710, NSUMA, FCUB, UPB, BUB, BUBS                  00004310
0311                  4710 FORMAT("0", "NSUMA, FCUB, UPB, BUB, BUBS ", 2I10, 3F15.4)                  00004320
C                  COMPARE UPPER BOUND WITH BEST UPPER BOUND                  00004330
C                  IF UPB IS LESS THAN BUB, SET IT AS BUB AND                  00004340
C                  NOTE THE SOLUTION                  00004350
0312                  4750 IF (UPB.GE.BUB) GO TO 5100                  00004360
0313                  4770 BUR=UPB                  00004370
0314                  BUBS= BUB/ (1.0+EPS)                  00004380
0315                  IBNOD=NOD                  00004385
0316                  DO 4800 J=1,N                  00004390
0317                  4800 BSOLX(J)=SOLX(J)                  00004400
0318                  DO 4850 K=1,P                  00004410
0319                  4850 BSOLY(K)=FUB(K)
C                  *****COMPARE LOWB WITH BUBS. IF LOWB IS GREATER                  00004430
C                  THAN OR EQUAL TO BUBS, THEN BACKTRACK                  00004440
0320                  4900 IF (LOWB.GE.BUBS)GO TO 6200                  00004450
C                  *****IF FIX(J) VALUES ARE 1 FOR EACH J, THEN BACKTRACK                  00004480
0321                  5100 IF (LQ1.EQ.N) GO TO 6200                  00004500
C                  *****APPLY THE BOUNDING RULE*****                  00004510
C                  IF THE DIFFERENCE BETWEEN C(I,J) AND MINC(J) IS                  00004515
C                  GREATER THAN THE DIFFERENCE BETWEEN BUBS AND                  00004520
C                  LOWB, THEN CX(I,J)=2                  00004525
C                  *****APPLY BRANCHING RULE AND FIND BRI, THE NEXT                  00004530
C                  VARIABLE FOR LEFT BRANCHING,                  00004540
C                  FIND NMINC(J), THE NEXT HIGHER VALUE THAN MINC(J)                  00004550
C                  AND DIFBR(J), THE DIFFERENCE BETWEEN THEM.                  00004555
0322                  DBOUND=BUBS-LOWB                  00004568
0323                  5200 DO 5250 J=1,N                  00004570
0324                  NMINC(J)=0.0                  00004580
0325                  DIFBR(J)=0.0                  00004590
0326                  5250 CONTINUE                  00004600
0327                  DO 5600 J=1,N                  00004610
C                  SKIP TO NEXT J IF FIX(J)=1                  00004620

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0328	IF (FIX(J).EQ.1) GO TO 5600		00004630
0329	LK=0		00004640
0330	I=1		00004650
C	SKIP C(I,J) IF CX(I,J)=2 & MOVE TO NEXT I		00004670
0331	5300 IF (CX(I,J).EQ.2) GO TO 5350		00004680
0332	IF (I.EQ.SOLX(J)) GO TO 5350		00004690
0333	IF ((C(I,J)-MINC(J)).GT.DBOUND) GO TO 5330		00004700
0334	IF ((I-LK).EQ.1) GO TO 5320		00004710
0335	IF (C(I,J).GE.NMINC(J)) GO TO 5400		00004720
0336	NMINC(J)=C(I,J)		00004730
0337	GO TO 5400		00004735
0338	5330 CX(I,J)=2		00004740
0339	NSX=NSX+1		00004742
0340	STX(NSX)=-(I*100+J)-1000000		00004745
0341	KT2(J)=KT2(J)+1		00004747
0342	IFI(KT2(J).LT.(M-1)) GO TO 5350		00004750
0343	INDI=SOLX(J)		00004752
0344	CX(INDI,J)=1		00004755
0345	NSX=NSX+1		00004758
0346	STX(NSX)=(INDI*100+J)+1000000		00004760
0347	FIX(J)=1		00004762
0348	LQ1=LQ1+1		00004764
0349	FIX(I,J)=INDI		00004766
0350	GO TO 5600		00004768
0351	5350 LK=LK+1		00004770
0352	5400 I=I+1		00004775
0353	IFI.LE.M) GO TO 5300		00004780
0354	5500 DIFBR(J)=NMINC(J)-MINC(J)		00004785
0355	5600 CONTINUE		00004790
0356	IF (ISTEP.NE.2) GO TO 5650		00004795
0357	DO 5620 I=1,M		00004820
0358	PRINT 290, I,(CX(I,J),J=1,N)		00004830
0359	5620 CONTINUE		00004850
0360	PRINT 297, (FIX(J),J=1,N)		00004860
C	IF FIX(J)=1 FOR ALL J, THEN BACKTRACK.		00004880
0361	5650 IF (LQ1.EQ.N) GO TO 6200		00004890
C	FIND MAXDIF, THE MAXIMUM DIFFERENCE DIFBR(J)		00004900
0362	LF=0		00004905
0363	DO 5800 J=1,N		00004910
0364	IF (FIX(J).EQ.1) GO TO 5690		00004915
0365	IF ((J-LF).EQ.1) GO TO 5660		00004920
0366	IF (DIFBR(J).LT.MAXDIF) GO TO 5800		00004925
0367	5660 MAXDIF=DIFBR(J)		00004930
0368	LJ=J		00004935
0369	GO TO 5800		00004940
0370	5690 LF=LF+1		00004943
0371	5800 CONTINUE		00004946
0372	IF (ISTEP.NE.2) GO TO 5840		00004950
0373	DO 5820 J=1,N		00004953
0374	IF (FIX(J).EQ.1) GO TO 5820		00004956
0375	PRINT 5810, J, NMINC(J), MINC(J), DIFBR(J)		00004960
0376	5810 FORMAT ('0',J,NMINC(J),MINC(J),DIFBR(J)', 15,3F15.4)		00004963
0377	5820 CONTINUE		00004966
C	*****BRANCHING VARIABLE BR1 CORRESPONDS TO MAXDIF*****		00004970
0378	5840 DO 5900 J=1,N		00004980
0379	IF (J.NE.LJ) GO TO 5900		00004990
0380	5850 BR1=SOLX(J)*100+J		00005000
0381	IF (ISTEP.EQ.0) GO TO 6020		00005010

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	MAIN	DATE = 80206	14/35/07
0382	PRINT 5880, BRI		00005020
0383	5880 FORMAT('0', BRI, I10)		00005030
0384	GO TO 6020		00005040
0385	5900 CONTINUE		00005050
	C *****UPDATE STX(INS) AND NSX*****		00005060
	C NSX REPRESENTS THE NUMBER OF VARIABLES IN STX(INS)		00005070
0386	6020 NSX=NSX+1		00005090
0387	6040 STX(NSX)=BRI		00005100
0388	IF (ISTEP.NE.2) GO TO 6100		00005150
0389	PRINT 6088, (STX(INS), INS=1,NSX)		00005160
0390	6088 FORMAT('0', STX(INS), I0I10, 1221/, I0X, I0I10))		00005170
	C *****MOVE TO THE NEXT (LEFT BRANCH) NODE AND APPLY		00005220
	C CAPACITY RULE		00005230
0391	6100 NOD=NOD+1		00005240
0392	6110 IF (ET.EQ.0.0) GO TO 6150		00005242*
0393	IF (INSET.EQ.1) GO TO 6147		00005244*
0394	IF (INET.EQ.1) GO TO 6150		00005246*
0395	CALL TIMET(INT)		00005248*
0396	ELTN=(INT-IT0)*26.04E-6		00005250*
0397	IF (ELTN.LT.ET) GO TO 6150		00005253*
0398	6120 PRINT 6125, NOD, ELTN, BUB, BUBS, IBNOD		00005256*
0399	6125 FORMAT ('0', 'WAS AT NODE', I6, ' AT ELAPSED TIME =', F10.4,		00005260*
	1 ' SECONDS.', /1X, ' BUB=', F15.4, ' BUBS=', F15.4,		00005263*
	2 ' AT NODE=', I7)		00005266*
0400	IBUB=BUB		00005267*
0401	IF (IBUB.EQ.9999999) GO TO 6146		00005268*
0402	6130 PRINT 6135, (BSOLX(J), J=1, N)		00005270*
0403	6135 FORMAT ('0', 'SOLUTION CORRESPONDING TO SUB IS', //1X,		00005273*
	1 '(BSOLX(J), J=1, N)', 10I8, 3I/16X, 10I8))		00005276*
0404	6140 PRINT 6145, (BSOLY(K), K=1, P)		00005280*
0405	6145 FORMAT (/1X, '(BSOLY(K), K=1, P)', 10I8, 2I/16X, 10I8))		00005290*
0406	6146 INET=1		00005292*
0407	INIS=ISTEP		00005294*
0408	INSET=1		00005296*
0409	ISTEP=2		00005298*
0410	GO TO 6150		00005300*
0411	6147 ISTEP=INIS		00005302*
0412	INSET=0		00005304*
0413	6150 GO TO 272		00005306
	C *****END IF AT THE ROOT NODE*****		00005308
0414	6200 IF (NSX.EQ.0) GO TO 8100		00005310
0415	6220 IF (IABS(STX(NSX)).GT.1000000) GO TO 6500		00005320
0416	6250 BRO=STX(NSX)		00005330
0417	6270 STX(NSX)=BRO-1000000		00005340
0418	IF (ISTEP.EQ.0) GO TO 6308		00005390
0419	PRINT 6305, BRO		00005400
0420	6305 FORMAT('0', 'BRO ', I10)		00005410
0421	6308 IF (ISTEP.NE.2) GO TO 6330		00005420
0422	PRINT 6088, (STX(INS), INS=1,NSX)		00005430
	C *****MOVE TO THE NEXT (RIGHT BRANCH) NODE AND APPLY		00005490
	C CAPACITY RULE		00005500
0423	6330 NOD=NOD+1		00005510
0424	6410 IF (ET.EQ.0.0) GO TO 6450		00005512*
0425	IF (INSET.EQ.1) GO TO 6445		00005516*
0426	IF (INET.EQ.1) GO TO 6450		00005518*
0427	CALL TIMET(INT)		00005520*
0428	ELTN=(INT-IT0)*26.04E-6		00005523*
0429	IF (ELTN.LT.ET) GO TO 6450		00005526*

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\*Relevant to the TIMET Subroutine

FORTRAN IV G LEVEL 21	MAIN	DATE = 80206	14/35/07
0430	6420 PRINT 6125, NOD, ELTN, BUB, BUBS, IBNOD		00005528*
0431	IBUB=BUB		00005530*
0432	IF (IBUB.EQ.9999999) GO TO 6442		00005532*
0433	6430 PRINT 6135, (BSOLX(J),J=1,N)		00005533*
0434	6440 PRINT 6145, (BSOLY(K), K=1,P)		00005536*
0435	6442 INET=1		00005538*
0436	INIS=ISTEP		00005540*
0437	INSET=1		00005542*
0438	ISTEP=2		00005544*
0439	GO TO 6450		00005546*
0440	6445 ISTEP=INIS		00005548*
0441	INSET=0		00005550*
0442	6450 GO TO 210		00005552
0443	6500 IF ( STX(NSX).GT.1000000) GO TO 6520		00005555
0444	LX=STX(NSX)-1000000		00005560
0445	IX=LX/100		00005570
0446	JX=LX-IX*100		00005580
0447	CX(IX,JX)=0		00005590
0448	KT2(JX)=KT2(JX)-1		00005595
0449	GO TO 6550		00005600
0450	6520 LX= STX(NSX)-1000000		00005610
0451	IX=LX/100		00005620
0452	JX=LX-IX*100		00005630
0453	CX(IX,JX)=0		00005640
0454	FIX(JX)=0		00005650
0455	LQ1=LQ1-1		00005660
0456	6550 NSX=NSX-1		00005670
0457	GO TO 6200		00005700
	C *****PRINT THE OUTPUT*****		00005730
0458	8100 IBUB=BUB		00005740
0459	CALL TIMET(IT1)		00005750*
0460	ELT1=(IT1-ITO)*26.04E-6		00005760*
0461	PRINT 8105, ELT1		00005770*
0462	8105 FORMAT ('0',//1X, 'ELAPSED TIME IN SECONDS=', F15.8)		00005780*
0463	PRINT 8120, NOD		00005790
0464	8120 FORMAT ('0', 'TOTAL NUMBER OF NODES EXPLORED =', I3)		00005800
0465	IF (IBUB.EQ.9999999) GO TO 8350		00005810
0466	8130 PRINT 8150		00005820
0467	8150 FORMAT ('0', 'NOTE: 1. FOLLOWING X(I,J) VARIABLES SHOW DESIGN',		00005830
	1 ' I TO WHICH ACTIVITY J IS ASSIGNED FOR J=1 TO N.',		00005840
	2 '/X, '2. IF EPSILON EPS WAS ASSIGNED A POSITIVE',		00005850
	3 ' (NON-ZERO) VALUE, THE SOLUTION MAY BE SUBOPTIMAL.',/)		00005860
0468	8180 PRINT 8200, (BSOLX(J),J=1,N)		00005870
0469	8200 FORMAT('0',T55, 'OPTIMAL SOLUTION',/1X,T55,		00005880
	1 '-----',//1X, 'X(I,J) WITH VALUE 1:',10I8,		00005890
	2 31/,21X,10I8))		00005900
0470	8220 PRINT 8250, (BSOLY(K), K=1,P)		00005910
0471	8250 FORMAT ('0', 'Y(K) VALUES:', 8X, 10I8, 21/,21X,10I8))		00005920
0472	8280 PRINT 8300,IBUB		00005930
0473	8300 FORMAT ('0', 'OPTIMAL VALUE OF OBJECTIVE FUNCTION:', I15//)		00005940
0474	GO TO 8500		00005950
0475	8350 PRINT 8400		00005960
0476	8400 FORMAT ('0', ' PROBLEM DOES NOT HAVE A FEASIBLE SOLUTION',		00005970
	1 '/1X, ' BECAUSE THE CAPACITY CONSTRAINTS CANNOT',		00005980
	2 '/1X, ' BE SATISFIED.',/)		00005990
0477	8500 PRINT 8550		00006000
0478	8550 FORMAT ('0', '*****NORMAL END OF JOB*****',/)		00006010
0479	8600 STOP		00006020
0480	END		00006030

\*Relevant to the TIMET Subroutine

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APPENDIX B

DIC<sup>I</sup>ONARY OF ZIPC<sup>A</sup>P SYM<sup>B</sup>OLIC NAMES

<u>SYMBOLIC NAME</u>	<u>DEFINITION</u>
A(I,J)	VARIABLE COST OF ACTIVITY J USING DESIGN I
B(K)	FIXED COST OF FACILITY K
BR0	BRANCHING VARIABLE WITH VALUE 0, I.E., RIGHT-BRANCH VARIABLE X(I,J)
BR1	BRANCHING VARIABLE WITH VALUE 1, I.E., LEFT-BRANCH VARIABLE X(I,J)
BSOLX(J)	CURRENT BEST SOLUTION CONTAINING X(I,J) VARIABLES
BSOLY(K)	CURRENT BEST SOLUTION CONTAINING Y(K) VARIABLES
BSUM	B(K) * D(I, J, K)/S(K) SUMMED OVER APPLICABLE K
BUB	BEST UPPER BOUND
BUBS	BEST UPPER BOUND IF SOLUTION MAY BE SUBOPTIMAL WITHIN EPS. BUBS = BUB/(1 + EPS)
C(I,J)	COST ELEMENTS FOR THE RELAXED PROBLEM
CX(I,J)	CONTAINS FIXED AND FREE X(I,J) VARIABLES. FIXED VARIABLES ARE ASSIGNED A VALUE OF 1 OR 2, AND FREE VARIABLES ARE ASSIGNED A VALUE 0. A VALUE OF 1 IMPLIES THAT THAT PARTICULAR X(I,J) HAS A VALUE 1, AND FIX(J) IS FIXED AS 1. A VALUE OF 2 IMPLIES THAT THAT PARTICULAR VARIABLE SHOULD NOT BE CONSIDERED FOR CURRENT COMPUTATIONS CX(I,J) IS UPDATED BY THE CAPACITY RULE, THE BOUNDING RULE, AND THE RULES FOR BRANCHING AND BACKTRACKING
D(I,J,K)	CAPACITY REQUIRED AT FACILITY K FOR ACTIVITY J WHEN ACTIVITY J USES DESIGN I
DBOUND	DIFFERENCE BETWEEN BOUNDS BUBS AND LOWB
DIFBR(J)	DIFFERENCE IN C(I,J) VALUES FOR SELECTING BRANCHING VARIABLE
E(I,K)	HAS VALUE 1 IF DESIGN I USES FACILITY K, OTHERWISE THE VALUE IS 0
ELT1	ELAPSED TIME IN SECONDS TO EXECUTE THE PROGRAM OBTAINED FROM THE TIMET SUBROUTINE

ELTN	ELAPSED TIME IN SECONDS TO COMPARE WITH ET IF SPECIFIED
EPS	EPSILON VALUE IF A SUBOPTIMAL SOLUTION ACCEPTABLE, E.G., EPS OF 0.002 IMPLIES THAT THE SOLUTION MAY BE SUBOPTIMAL, BUT IS GUARANTEED TO HAVE A VALUE WITHIN 0.2% OF THE OPTIMAL VALUE. EPS IS 0 IF AN OPTIMAL SOLUTION IS DESIRED
ET	ELAPSED TIME, IF SPECIFIED, AT WHICH THE NODE AND BOUNDS INFORMATION IS PRINTED
FC	FIXED COST FOR COMPUTING LOWER BOUND
FCUB	FIXED COST FOR COMPUTING UPPER BOUND
FIX(J)	VECTOR OF FIXED COLUMNS. A COLUMN IS FIXED IF IT HAS A FIXED X(I,J) VARIABLE WITH VALUE 1
FIXI(J)	VECTOR SPECIFYING INDEX I CORRESPONDING TO FIX(J) OF VALUE 1
FLB(K)	VECTOR OF FACILITIES FOR COMPUTING LOWER BOUND
FUB(K)	VECTOR OF FACILITIES FOR COMPUTING UPPER BOUND
I	INDEX FOR DESIGNS
IBALD	BALANCE AVAILABLE CAPACITY FOR A GIVEN FACILITY
IBNOD	NODE CORRESPONDING TO THE BEST UPPER BOUND
IBUB	BEST UPPER BOUND
ICAPR	OPTION TO USE CAPACITY RULE
IDIFD	DIFFERENCE BETWEEN A D(I,J,K) VALUE AND MIND(J)
IINPT	OPTION TO LIST INPUT DATA
IMIN	MINIMUM OVER INDEX I FOR A GIVEN COLUMN J
INET	INDICATOR FOR ET
INIS	HAS THE SAME VALUE AS ISTEP
INS	INDEX FOR STACK OF VARIABLES STX
INSET	INDEX FOR DISPLAYING DETAILED STEPS IF OPTION ET IS SPECIFIED
INT	INTERMEDIATE CALL TO TIMET SUBROUTINE

ISTEP            OPTION TO LIST SUMMARY OF INTERMEDIATE STEPS, DETAILED  
 STEPS, OR TO SKIP THE LISTING  
  
 IT0            INITIAL CALL TO TIMET SUBROUTINE  
  
 IT1            FINAL CALL TO TIMET SUBROUTINE  
  
 IUNCAP        OPTION TO INDICATE IF THE PROBLEM BEING SOLVED IS  
 CAPACITATED OR UNCAPACITATED  
  
 IX            INDEX I OF VARIABLE X(I,J)  
  
 J            INDEX FOR ACTIVITIES  
  
 JX            INDEX J OF VARIABLE X(I,J)  
  
 K            INDEX FOR FACILITIES  
  
 KT2(J)        A COUNTER FOR CX(I,J) VALUES SET EQUAL TO 2 FOR  
 COLUMN J  
  
 LK            A COUNTER  
  
 LOWB        LOWER BOUND  
  
 LQ1            A COUNTER FOR CX(I,J) VALUES SET EQUAL TO 1  
  
 LQ2            A COUNTER FOR CX(I,J) VALUES SET EQUAL TO 2 IN A  
 CYCLE IN WHICH ALL THE FACILITIES ARE EXAMINED  
  
 LR2            A COUNTER SET EQUAL TO LQ2 AT THE END OF A CYCLE WHEN  
 APPLYING THE CAPACITY RULE  
  
 P            NUMBER OF FACILITIES  
  
 S(K)        TOTAL CAPACITY OF FACILITY K  
  
 SOLX(J)      SOLUTION SHOWING INDEX 1 OF X(I,J) VARIABLES  
 CORRESPONDING TO J=1, 2, ..., N  
  
 STX(INS)     STACK OF FIXED X(I,J) VARIABLES FOR BRANCHING AND BACK-  
 TRACKING. AN X(I,J) RECORDED IN CX(I,J) AS 1 DUE  
 TO THE BRANCHING RULE IS RECORDED IN STX(INS) AS  
 $X * 100 + J$ . AN X(I,J) RECORDED IN CX(I,J) AS  
 1 DUE TO THE CAPACITY RULE OR THE BOUNDING RULE IS  
 RECORDED IN STX(INS) AS  $(X * 100 + J) + 1,000,000$ .  
 AN X(I,J) RECORDED IN CX(I,J) AS 2 IS RECORDED  
 IN STX(INS) AS  $-(X * 100 + J) - 1,000,000$   
  
 TIMET        A SUBROUTINE FROM THE PLI LIBRARY TO RECORD THE  
 EXECUTION TIME

UPB

UPPER BOUND

ZIPCAP

NAME OF THE PROGRAM, AN ACRONYM FOR ZERO-ONE INTEGER  
PROGRAM TO SOLVE MULTIACTIVITY MULTIFACILITY CAPACITY-  
CONSTRAINED ASSIGNMENT PROBLEMS

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APPENDIX C

ANNOTATED OUTPUT FOR TEST PROBLEM 1

OPTIONS SELECTED : INPUT=1 ICAP=1 ISTEP=1 IUNCAP=0 EPS= 0.0 ET= 0.0

INPUT DATA

NUMBER OF DESIGNS (M)= 3  
NUMBER OF ACTIVITIES (N)= 4  
NUMBER OF FACILITIES (P)= 5

VARIABLE COST MATRIX A(I,J)

I= 1	9847	15716	4450	8925
I= 2	10192	10542	4422	8118
I= 3	11019	9750	4409	8337

FIXED COST VECTOR B(K)

I= 1	1750	2000	1750	1350	1000
------	------	------	------	------	------

CAPACITY LIMIT VECTOR S(K)

K= 1	400	400	1000	400	400
------	-----	-----	------	-----	-----

CAPACITY USAGE MATRIX D(I,J,K)

K= 1	I= 1	0	0	0	0
K= 2	I= 1	180	60	261	149
K= 3	I= 1	154	36	175	100

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(C)

(D)

	K= 2	K= 3	K= 4	K= 5
I= 1	0	525	0	0
I= 2	0	186	0	0
I= 3	154	154	36	36
			1011	1011
			144	144
			60	60
			36	36
			175	175
			100	100

	K= 2	K= 3	K= 4	K= 5
I= 1	0	0	0	0
I= 2	0	0	0	0
I= 3	154	154	36	36
			1011	1011
			144	144
			60	60
			36	36
			175	175
			100	100

DESIGN-FACILITY MATRIX E(I,I,K)

	K= 2	K= 3	K= 4	K= 5
I= 1	0	0	0	0
I= 2	1	1	1	1
I= 3	1	1	1	1

## SUMMARY OF STEPS

	MINS C. FC. LOWs	37229.3594	0	37229.3594
BRI	101			

NODE NUMBER	2			
MNSC, FC, LNSB	35926.3633			
BRI	203			
NODE NUMBER	3			
MNSC, FC, LNSB	32356.0000			
NSUMA, FCUB, UPB, BLB, SUMS	32356			
BRI	203			
NODE NUMBER	4			
MNSC, FC, LNSB	32324.0000			
NSUMA, FCUB, UPB, BLB, SUMS	32324			
BRI	101			
NODE NUMBER	5			
MNSC, FC, LNSB	33476.0000			
NSUMA, FCUB, UPB, BLB, SUMS	33476			
BRI	203			

ELAPSED TIME IN SECONDS = 0.06226163

TOTAL NUMBER OF NODES EXPLORIED = 5

NOTE: 1. FOLLOWING X(i,j) VARIABLES SHOW DESIGN I TO WHICH ACTIVITY j IS ASSIGNED FOR j=1 TO N.  
 2. IF EPSILON EPS WAS ASSIGNED A POSITIVE (NON-ZERO) VALUE, THE SOLUTION MAY BE SUBOPTIMAL.

OPTIMAL SOLUTION

X(i,j) WITH VALUE 1:	1	3	3	2
Y(i) VALUES:	1	1	1	1
OPTIMAL VALUE OF OBJECTIVE FUNCTION:	40174			

\*\*\*\*\*NORMAL END OF JOB\*\*\*\*\*

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APPENDIX D

ANNOTATED OUTPUT FOR TEST PROBLEM 2

OPTIONS SELECTED : LIMIT=1 ICAPR=1 ISTEP=0 IUNCAP=0 EPS= 0.00200 ET= 0.0

INPUT DATA

NUMBER OF DESIGNS (M)= 5  
NUMBER OF ACTIVITIES (N)= 4  
NUMBER OF FACILITIES (P)= 8

VARIABLE COST MATRIX A(I,J)

I= 1	194951	218871	155094	104056
I= 2	235277	272087	143264	136641
I= 3	196138	220718	160399	107481
I= 4	198751	224042	167046	112445
I= 5	190873	229169	128361	112498

FIXED COST VECTOR B(M)

14000	14000	14000	14000	14000
				31000

CAPACITY LIMIT VECTOR S(K)

350	350	350	200	200
			700	700

CAPACITY USAGE MATRIX U(I,J,K)

K= 1

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(c)

I= 1	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100
I= 2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
I= 3	80	90	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100
I= 4	40	30	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50
I= 5	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
K= 2	I= 1	80	90	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100
I= 2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
I= 3	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60
I= 4	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40
I= 5	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
K= 3	I= 1	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60
I= 2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
I= 3	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60
I= 4	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40
I= 5	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
K= 5	I= 1	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60
I= 2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
I= 3	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60	60
I= 4	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40	40
I= 5	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

(c)

I= 4	40	30	50	120
I= 5	80	90	100	180
K= 6				
I= 1	240	180	300	450
I= 2	0	0	0	0
I= 3	120	90	150	360
I= 4	0	0	0	0
I= 5	30	50	120	

K= 7				
I= 1	180	180	200	360
I= 2	0	0	0	0
I= 3	60	60	100	240
I= 4	0	0	0	0
I= 5	180	180	200	360
K= 8				
I= 1	200	150	250	600
I= 2	40	30	50	180
I= 3	0	0	0	0
I= 4	0	0	0	0
I= 5	120	90	150	300

DESIGN-FACILITY MATRIX E(I,J,K)

I= 1	1	1	1	1	1	1	1
I= 2	0	0	0	0	0	0	0
I= 3	1	1	1	1	1	1	1
I= 4	0	0	0	0	0	0	0
I= 5	1	1	1	1	1	1	1

1  
0  
5  
0  
0  
0  
1  
1  
1  
1  
1

ELAPSED TIME IN SECONDS =

0.09392625

TOTAL NUMBER OF NODES EXPLORED = 11

NOTE: 1. FOLLOWING X(I,J) VARIABLES SHOW DESIGN I TO WHICH ACTIVITY J IS ASSIGNED FOR J=1 TO N.  
 2. IF EPSILON EPS WAS ASSIGNED A POSITIVE (NON-ZERO) VALUE, THE SOLUTION MAY BE SUBOPTIMAL.

OPTIMAL SOLUTION

X(I,J) WITH VALUE 1:  
 Y(I,J) VALUES:  
 1 1 1 1 1 1 1 1 1 1

OPTIMAL VALUE OF OBJECTIVE FUNCTION: 779502

\*\*\*\*\*NORMAL END OF JOB\*\*\*\*\*

(e)  
(f)

(g)

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APPENDIX E

ANNOTATED OUTPUT FOR TEST PROBLEM 3

OPTIONS SELECTED : LINP=1 ICAP=0 ISSTEP=0 JUNCAP=1 EPS=0.0 ET= 30.000

INPUT DATA

NUMBER OF DESIGNS (M)= 10  
 NUMBER OF ACTIVITIES (N)= 8  
 NUMBER OF FACILITIES (P)= 8

VARIABLE COST MATRIX A(I,J)

I= 1	21	13	20	8	15	10	35	20
I= 2	26	13	27	14	17	13	40	25
I= 3	25	12	26	13	16	12	20	23
I= 4	23	9	22	20	19	20	50	19
I= 5	23	20	24	13	20	15	40	17
I= 6	23	20	24	13	20	15	40	17
I= 7	26	13	26	20	15	15	35	20
I= 8	30	20	40	23	30	20	55	25
I= 9	26	13	26	20	20	15	40	15
I= 10	26	46	44	50	19	20	40	20

FIXED COST VECTOR B(I)

10	8	4	10	14	30	20	50
----	---	---	----	----	----	----	----

(c)

CAPACITY LIMIT VECTOR SICK

CAPACITY USAGE MATRIX OUT, J, K)

IN 1	IN 1	IN 2	IN 3	IN 4	IN 5	IN 6	IN 7	IN 8	IN 9	IN 10	IN 2	IN 1	IN 2	IN 3	IN 4
------	------	------	------	------	------	------	------	------	------	-------	------	------	------	------	------

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6

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6

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(c)

- - - - - 1 0 0 1 0 - - - - - 0 0 0 1 0

- - - - - 1 0 0 1 0 - - - - - 1 0 0 1 0

- - - - - 1 0 0 1 0 - - - - - 1 0 0 1 0

- - - - - 1 0 0 1 0 - - - - - 1 0 0 1 0

- - - - - 1 0 0 1 0 - - - - - 1 0 0 1 0

- - - - - 1 0 0 1 0 - - - - - 1 0 0 1 0

- - - - - 1 0 0 1 0 - - - - - 1 0 0 1 0

- - - - - 1 0 0 1 0 - - - - - 1 0 0 1 0

2 3 4 5 6 7 8 9 10 7 8 9 10 5 6 7 8 9 10 8

2

DESIGN-FACILITY MATRIX ELEMENT

ELAPSED TIME IN SECONDS: 0.57446837  
 TOTAL NUMBER OF MODES EXPLORED = 35

NOTE: 1. FOLLOWING X(I,J) VARIABLES SHOW DESIGN I TO WHICH ACTIVITY J IS ASSIGNED FOR J=1 TO N.  
 2. IF EPSILON EPS WAS ASSIGNED A POSITIVE (NON-ZERO) VALUE, THE SOLUTION MAY BE SUBOPTIMAL.

OPTIMAL SOLUTION

X(I,J) WITH VALUF I:	6	3	6	3	3	3	3	6
VALUFS:	1	1	1	0	1	1	0	1
OPTIMAL VALUE OF OBJECTIVE FUNCTION:	253							

\*\*\*\*\*NORMAL END OF JOBS\*\*\*\*\*

(e)  
(f)

(g)

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A VAULT FOR THE FUTURE  
IN THE YEAR 2036

THE STORY OF ENGINEERING IN THIS YEAR OF THE PLACING OF THE VAULT AND  
ENGINEERING HOPES FOR THE TOMORROWS AS WRITTEN IN THE RECORDS OF THE  
FOLLOWING GOVERNMENTAL AND PROFESSIONAL ENGINEERING ORGANIZATIONS AND  
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THE PURPOSE OF THIS VAULT IS INSPIRED BY AND IS DEDICATED TO  
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COMMUNITY, TO HIS NATION, AND TO OTHER NATIONS.

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